

Fair Scheduling for Wireless Mesh Networks with Multiple GWs

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Outline of Presentation

- ▶ Introduction
- ▶ Background & Related Work
- ▶ Limitations of Existing Work
- ▶ Proposed Approach
- ▶ Detailed Description of Fair Scheduling
- ▶ Performance Evaluation
- ▶ Conclusions & Future Work

Introduction

- ▶ **Wireless Mesh Network (WMN):**
 - Mesh Routers (MR)s: Often fixed, more resources (CPU, Memory) than traditional ad hoc nodes
 - Gateways (GW)s: Source and Destination of all Internet Traffic
 - Mesh Clients (MC)s: User devices such as laptops, PDAs, mobile phones etc.

Introduction: WMN

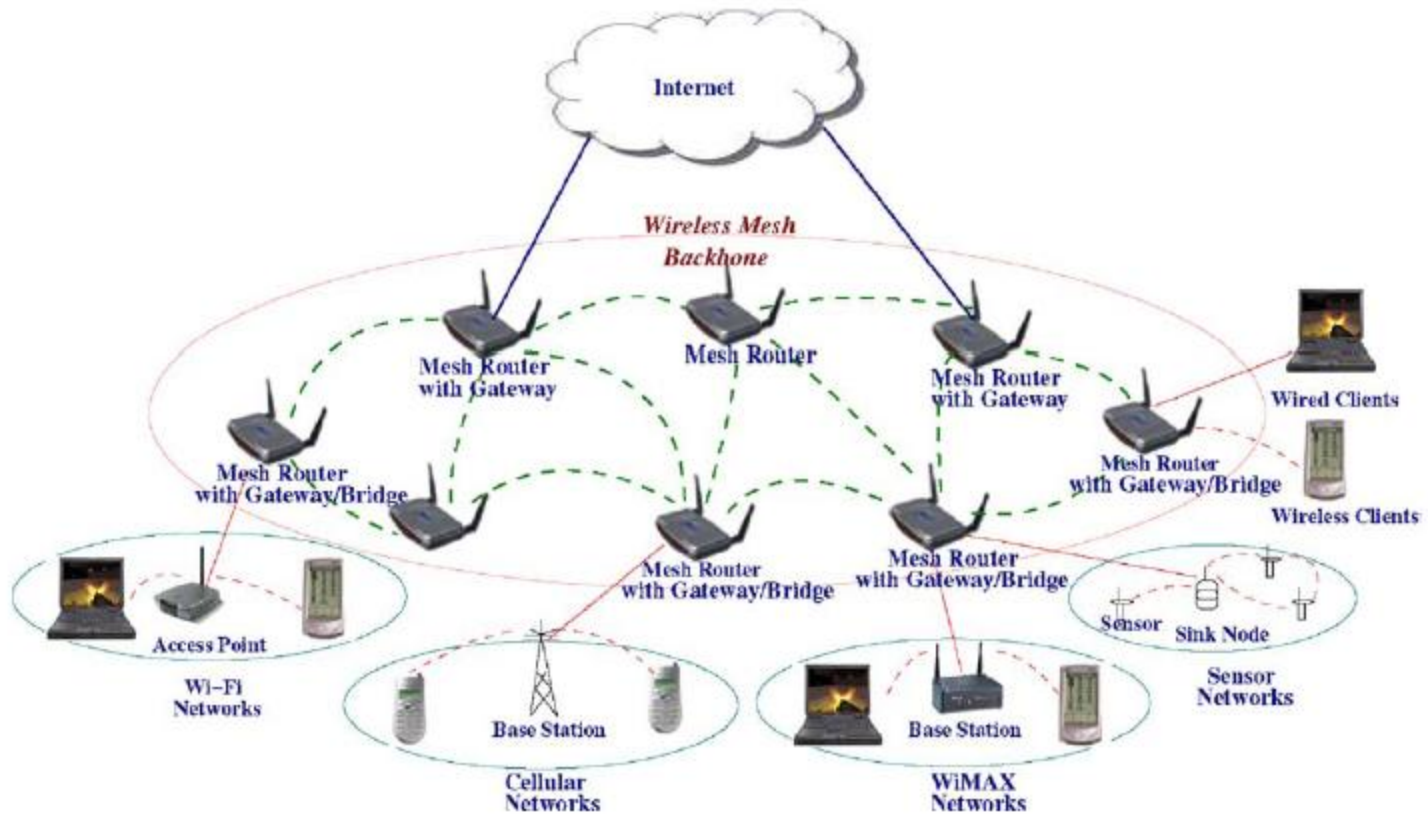


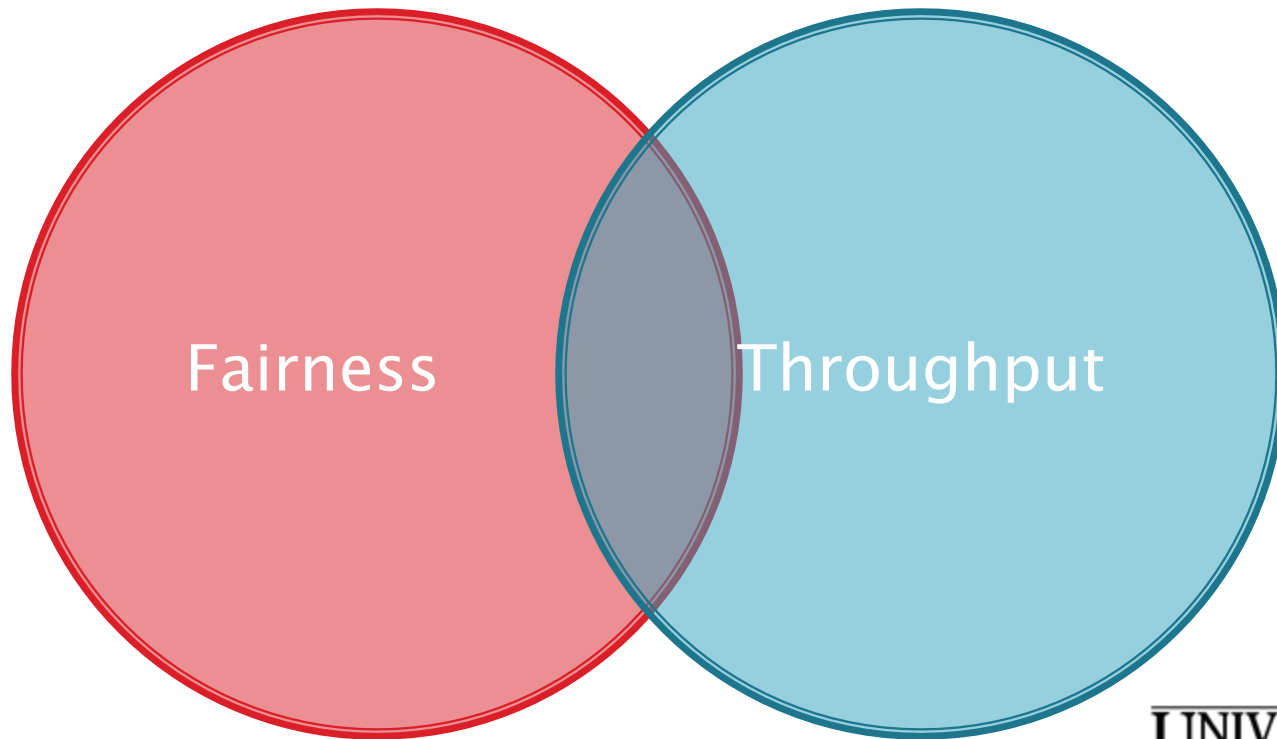
Fig. 3. Infrastructure/backbone WMNs.

[Akyildiz et al]

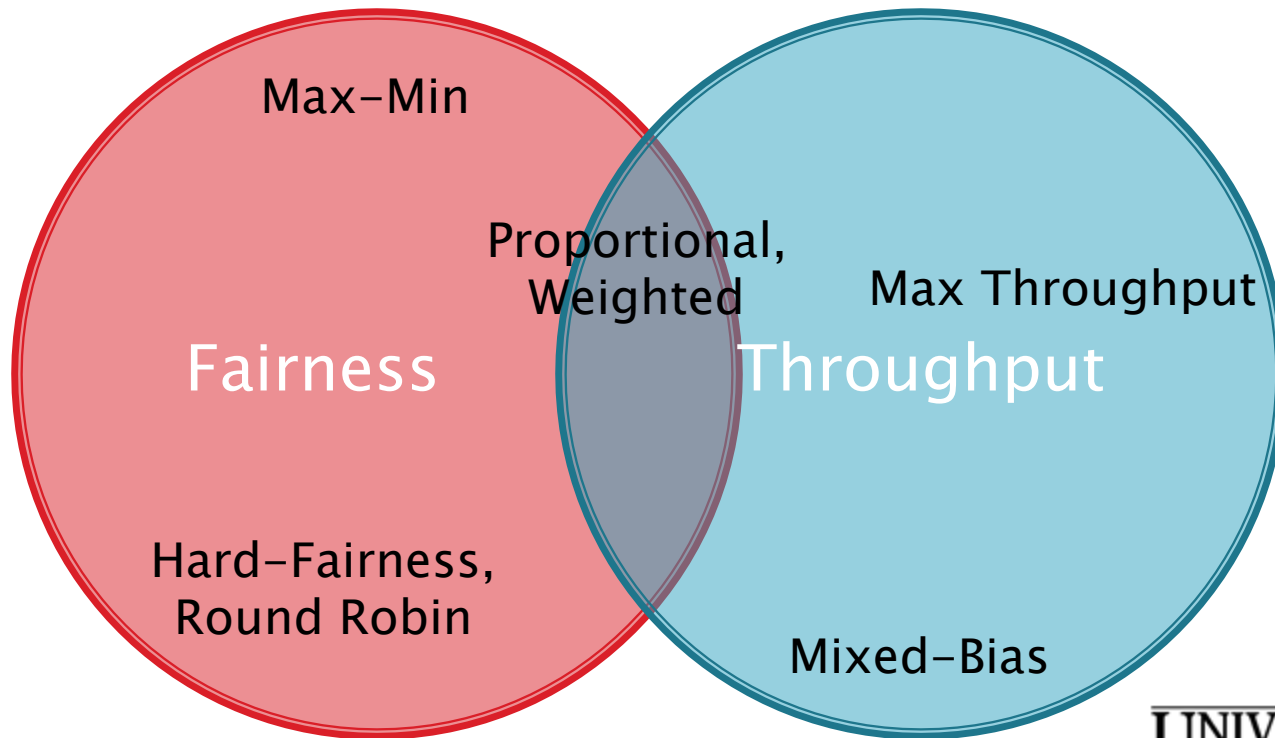
Introduction: Mobile WMN



Introduction: Definition of Fairness



Introduction: Definition of Fairness



Introduction: Motivation for Scheduling

- ▶ Equal Price, Equal Quality of Service (QoS)
- ▶ Avoid node starvation
- ▶ Wasted Resources

Background & Related Work

- ▶ Operating Systems
 - User and process scheduling in interactive OS's started in the 1960s and 1970s (multics, unix)
- ▶ Wired Networks, Wireless LAN (single hop)
- ▶ Ad hoc networks
- ▶ Other distributed computing
 - SHARCNET, World Community Grid

Background & Related Work

- ▶ Hard Fairness [Cao et al, Nelson et al]

$$\text{Resources assigned } (R) = \frac{1}{\# \text{ flows through the node}}$$

- ▶ Equal resources assigned to every node, flow, mesh router etc.
- ▶ Does not work well if some nodes do not require resources

Background & Related Work

- ▶ Max–min [Bejerano et al]
- ▶ Maximize the minimum amount of resources assigned to each node, or revert to an equal split
- ▶ Does not work well if there are large differences in requirements from node to node

Background & Related Work

- ▶ Proportional Fairness [Erwu et al, Gupta et al]

$$R = \frac{1}{c^\beta}$$

Where:

R is the resources allocated to the node

c is the characteristic which priority is given to, $c > 0$

β is the proportionality factor, $\beta > 0$

- ▶ Works well, but what if we would like to bias more strongly against certain characteristics while still retaining a minimum service level

Background & Related Work

- ▶ Mixed-Bias [Singh et al]

$$R = \frac{\alpha}{c^{\beta_1}} + \frac{1-\alpha}{c^{\beta_2}}$$

Where:

R is the resources allocated to the node

c is the characteristic which priority is given to, $c > 0$

β_1, β_2 are the proportionality factors, $\beta_1, \beta_2 > 0$

- ▶ Allows multiple strong and weak biasing against one characteristic

Limitations of Existing Work

▶ Fair Scheduling

- In [Salem et al] only scheduling proposed
 - No mechanism for maintaining and collecting the requirements
 - Performance evaluation based on length of schedule, not other metrics such as PDR or delay
 - Single GW

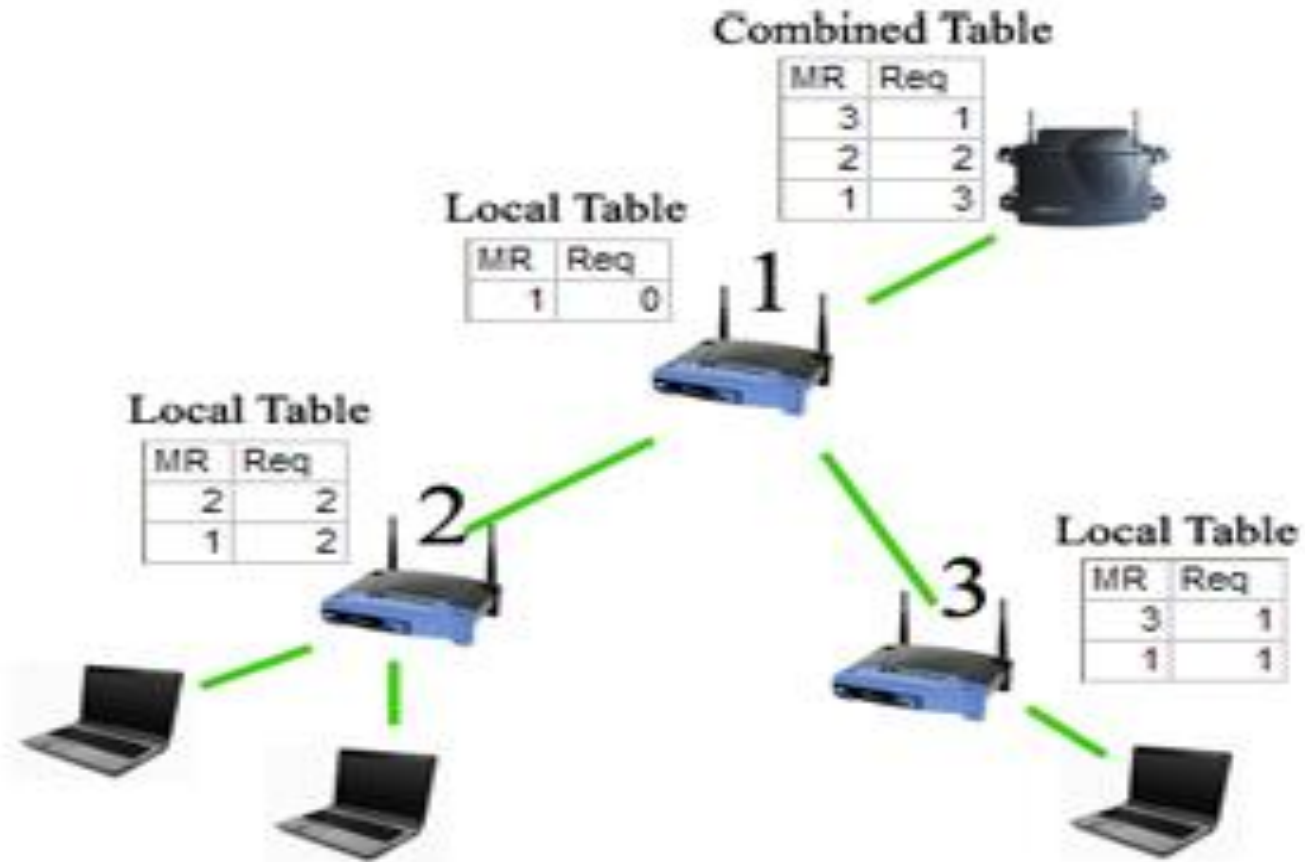
Proposed Approach

- ▶ Fair Scheduling for WMNs with Multiple Gateways
 - Distributed Requirement / Routing Tables
 - Requirement Propagation
 - Clique Generation / Compatibility Matrix
 - Implementation of Schedule Generation

Fair Scheduling for WMNs with Multiple Gateways

- ▶ Based on [Salem et al]
 - Our approach gives a mechanism for collecting the requirements and distributed the scheduling
 - We also relax the assumption of a single gateway
 - We evaluate the solution with a network simulation. The original authors evaluate based solely on scheduling lengths

Fair Scheduling



Distributed Routing & Requirement Tables 

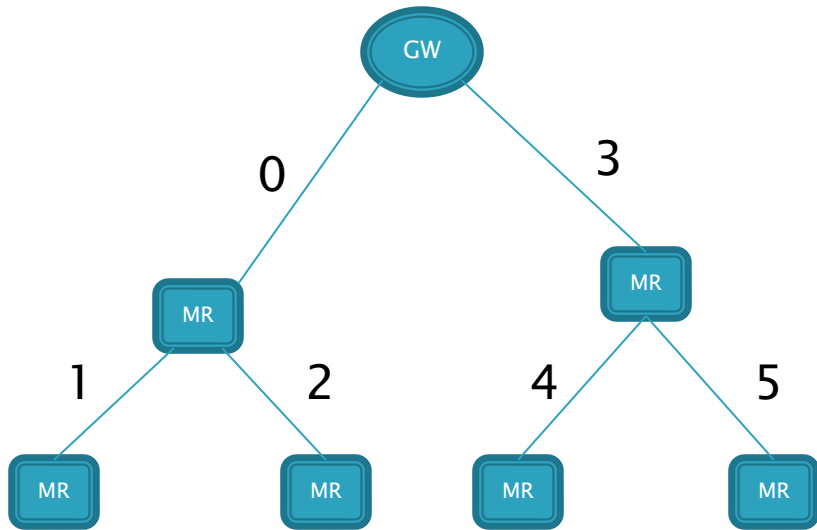
Fair Scheduling

Requirement Propagation Algorithm Pseudocode

1. *Associate MC with MR*
2. *Generate a Client Requirement at MR for the MC*
3. *For each link between MR and GW*
 - *Requirement(current-link) ++*
4. *For each Hop*
 - *Requirement(current-link) --*
5. *On Drop: For each link between MR and GW*
 - *Requirement(current-link) --*

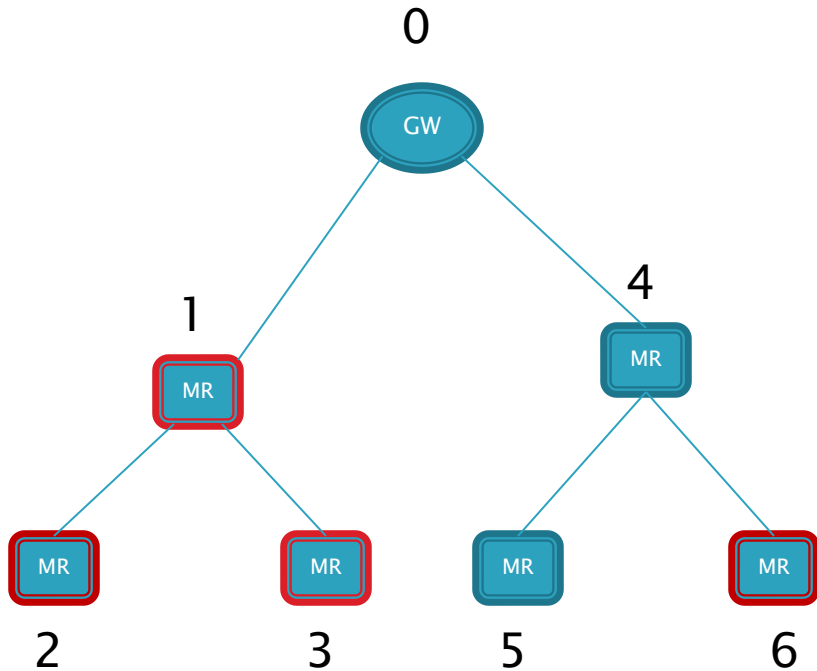
Algorithm 1 Requirement Propagation

Fair Scheduling: Clique Generation / Compatibility Matrix

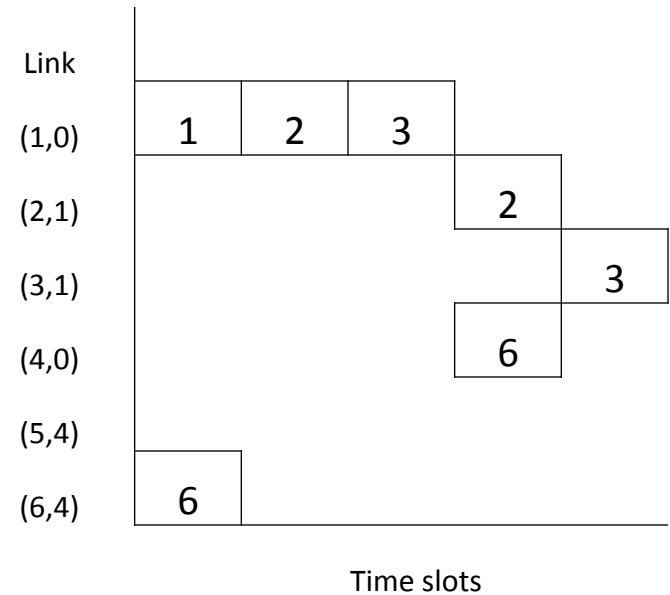


	0	1	2	3	4	5
0	1	0	0	0	1	1
1	0	1	0	1	1	1
2	0	0	1	1	1	1
3	0	1	1	1	0	0
4	1	1	1	0	1	0
5	1	1	1	0	0	1

Fair Scheduling: Schedule Generation



- Red outline denotes requirements at these MRs



Example STDMA Scheduling

Performance Evaluation

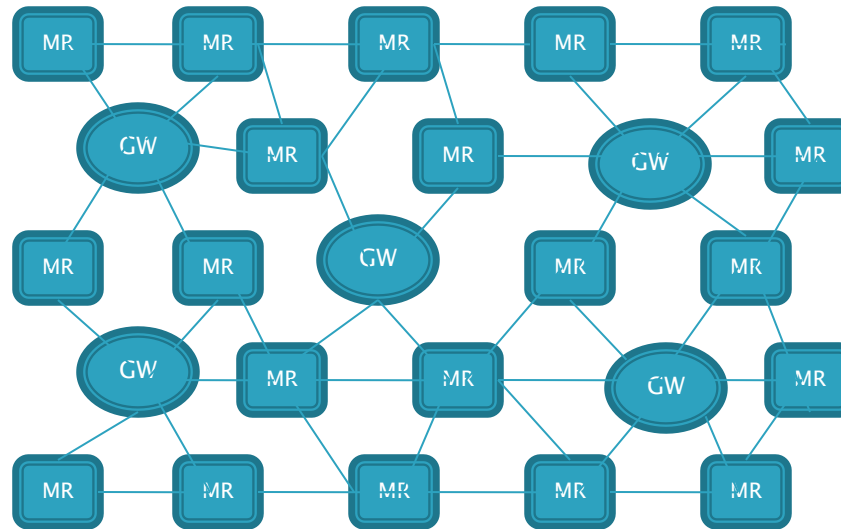
- ▶ Fair Scheduling for WMNs with Multiple GWs
 - Compared against single GW case
 - Compared against no fair scheduling case

Parameter	Value
Environment Dimensions	1000m x 1000m
Node Range	250m
Number of Mesh Routers	10 to 50
Number of Mesh Clients	250
Number of Gateways	1 to 5
Mean Packet Arrival	0.01 s
Mean Hop Delay	0.01 s
Retry Threshold	0.01 s

- Each run repeated 10 times for consistency

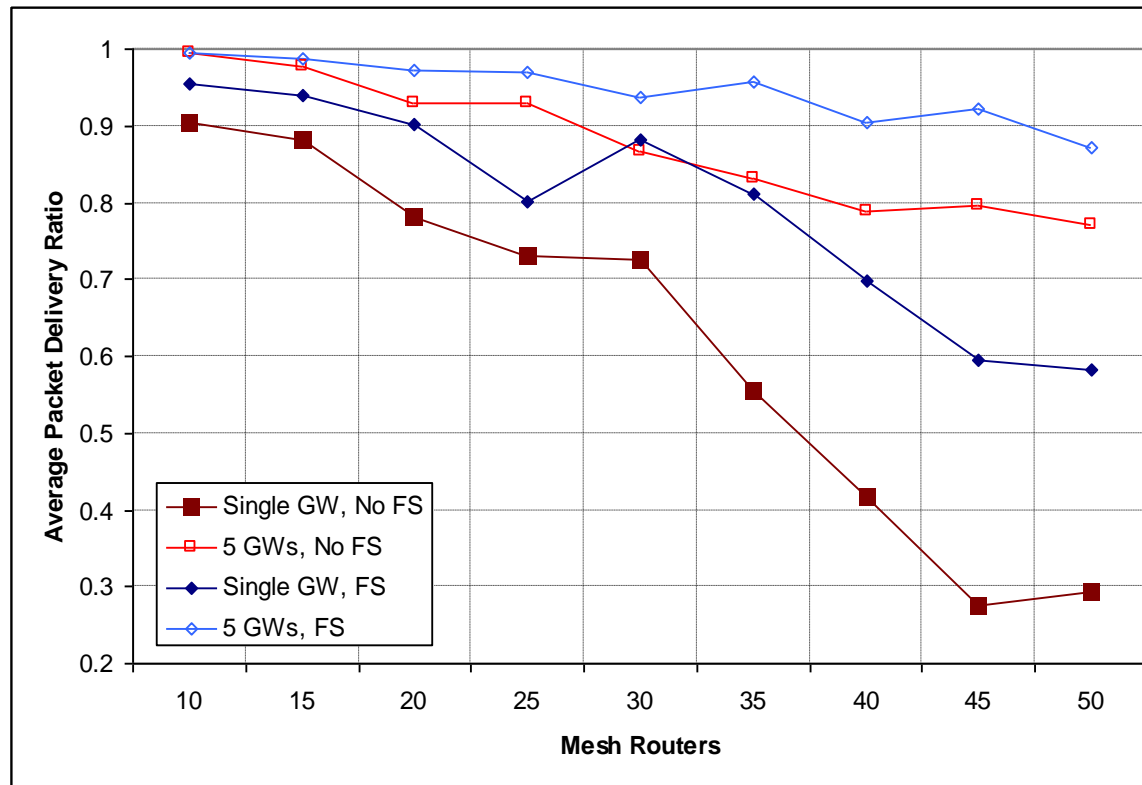
Traffic Model and Node Placement

- ▶ MRs arranged in a staggered grid topology so that each node can have at max 6 neighbours
- ▶ GWs spread out as much as possible
- ▶ Traffic routes towards closest GW
- ▶ Packet arrivals modeled with a Poisson process



Performance Evaluation

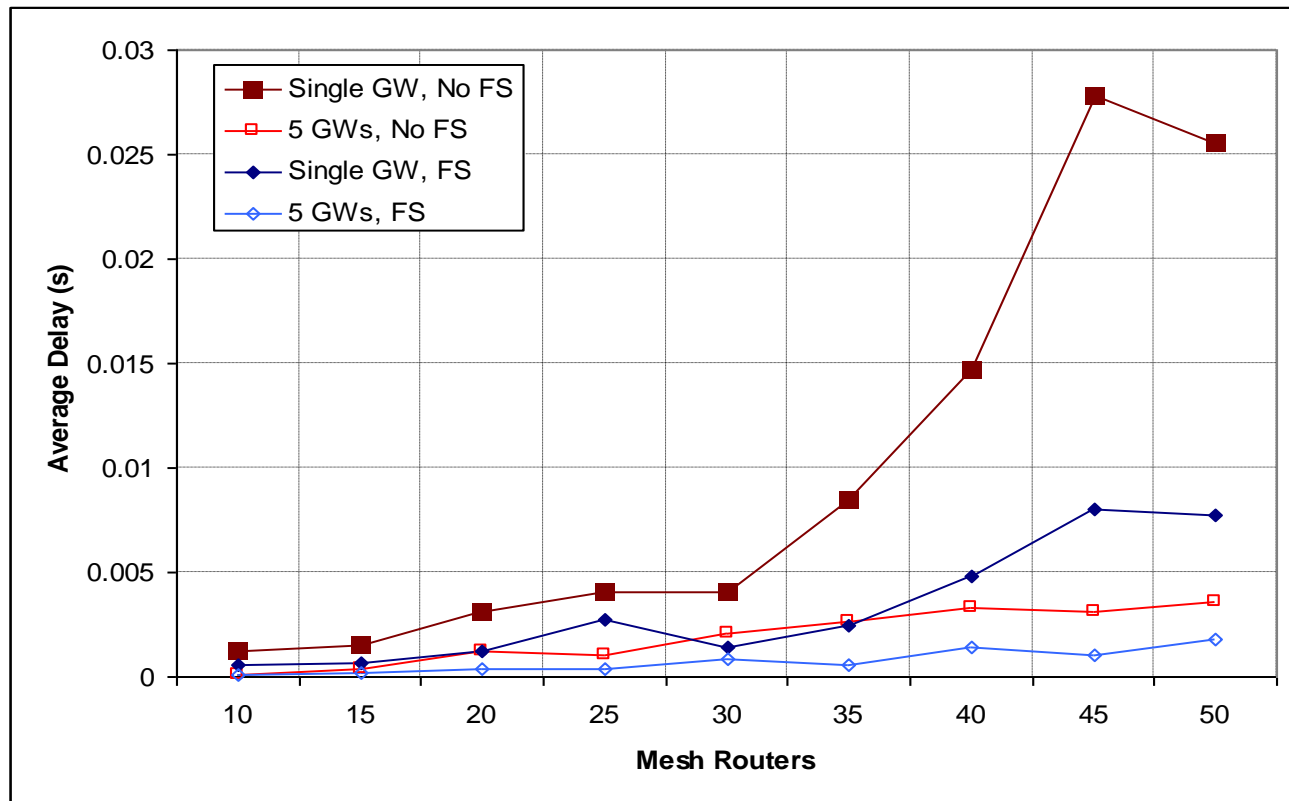
▶ Effect of Varying Mesh Routers on PDR



▶ Single GW performs poorly compared to Multi-GW

Performance Evaluation

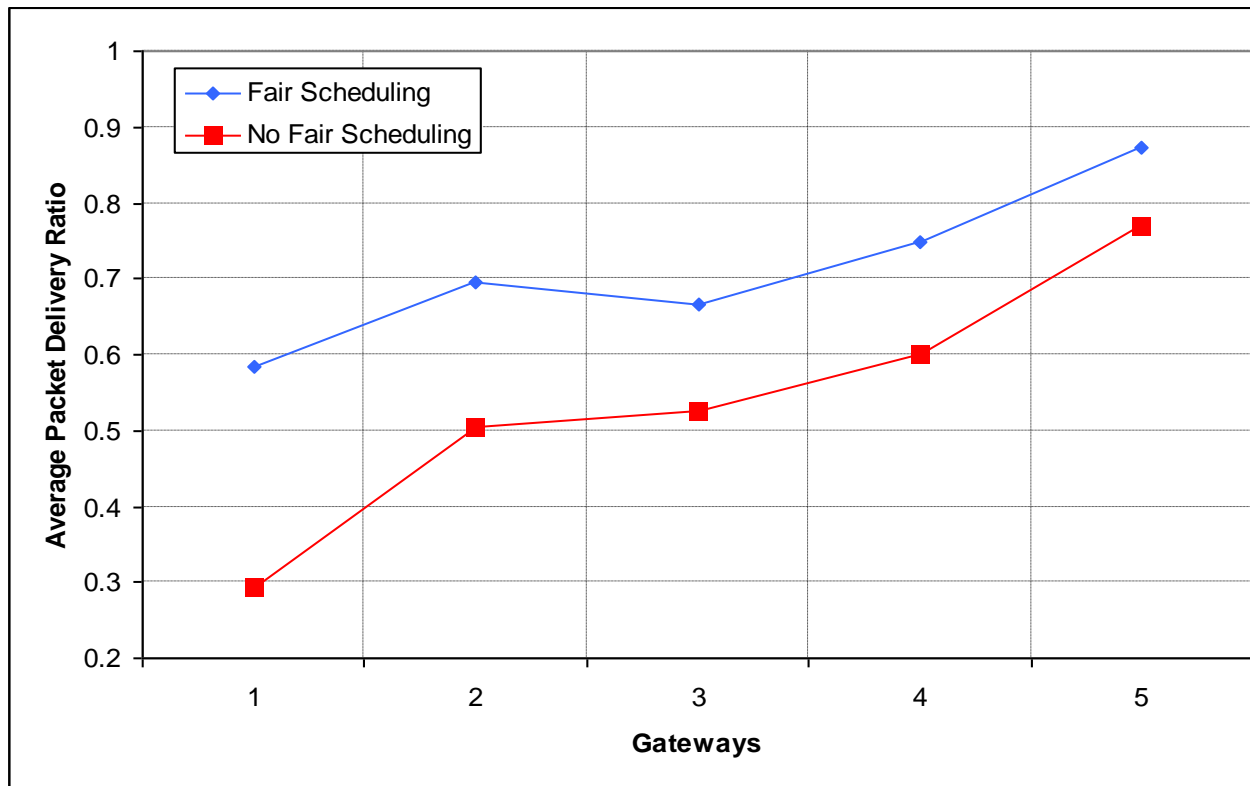
▶ Effect of Varying Mesh Routers on Delay



▶ High delay results in fewer successful packet arrivals

Performance Evaluation

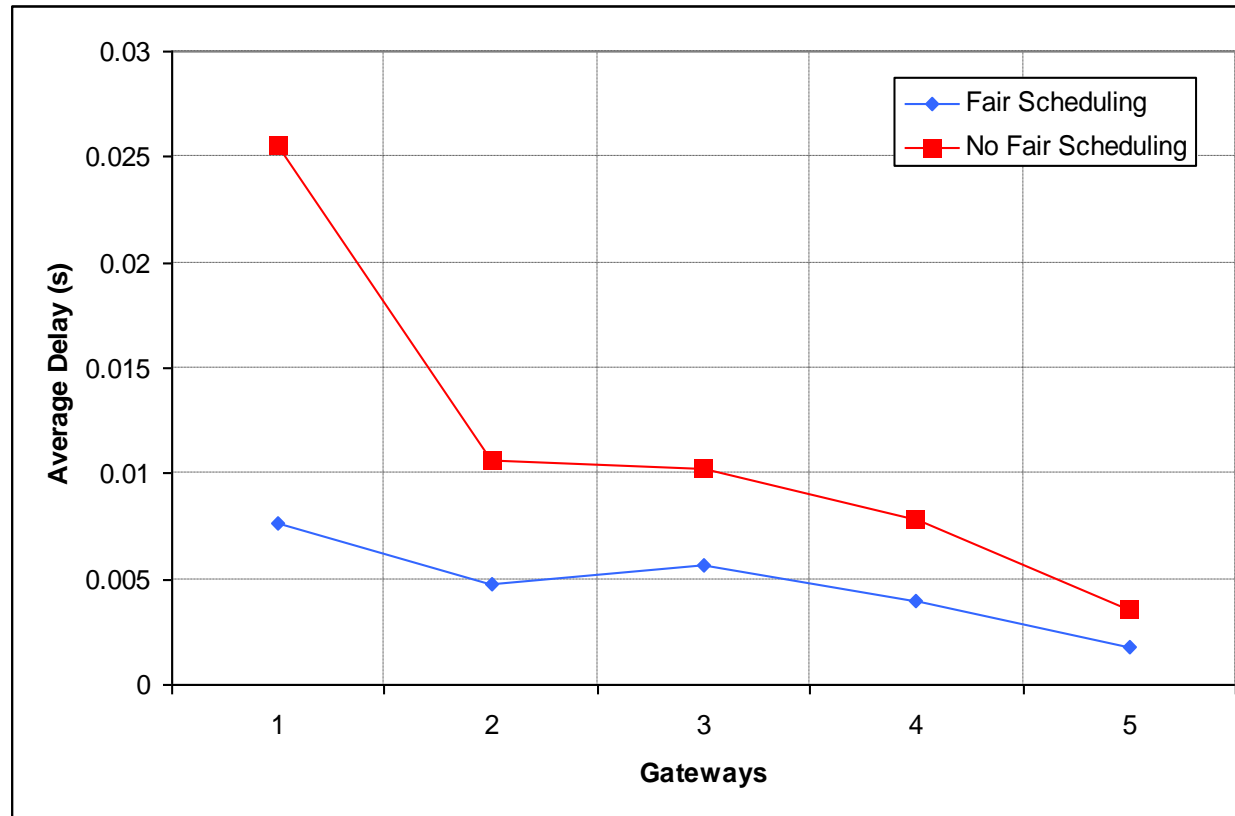
▶ Effect of Varying Gateways on PDR (50 MR)



- ▶ Multiple GWs allow increased performance most of the time
 - ▶ FS performs better than no FS

Performance Evaluation

- ▶ Effect of Varying Gateways on Delay (50 MR)



- ▶ Delay decreases significantly as more GWs are added since on average fewer hops are required

Conclusions & Future Work

- ▶ Fair Scheduling with for WMNs with multiple Gateways was proposed
- ▶ The main features of the proposed scheme are:
 - Distributed Requirement / Routing Tables
 - Requirement Propagation
 - Multiple GWs
- ▶ Simulation results with FS are improvement over no FS
- ▶ Multiple GWs are beneficial
- ▶ As a future work we will investigate:
 - The effect of mobility of mesh routers
 - Cross-layer design optimizations

Thank you

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